

# Invasion percolation in 2D

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Joint work with M. Damron and A. Sapozhnikov

# Bernoulli bond percolation

- Consider the square lattice  $\mathbb{Z}^2$  with its set of nearest neighbour bonds  $\mathbb{E}^2$ .
- Assign *uniform* $[0,1]$  random variables to each edge independently, denoted by  $\tau(e)$  for an edge  $e$ .
- For a given  $p$  we say that an edge  $e$  is  $p$ -open if  $\tau(e) \leq p$ .
- This model is called Bernoulli percolation with parameter  $p$ .
- The same model can be obtained by declaring each edge open with probability  $p$  and closed otherwise, independently of each other.

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# Some important results for Bernoulli percolation

- Let  $\Theta(p) = \mathbb{P}_p(0 \leftrightarrow \infty)$  be the percolation function.
- Let  $p_c = \inf\{0 \leq p \leq 1 : \Theta(p) > 0\}$  be the critical probability.
- For all  $p > p_c$  there is a unique infinite  $p$ -open cluster, denoted by  $\mathcal{C}_p$ .
- $p_c = 1/2$  and  $\Theta(p_c) = 0$ .
- Russo, Seymour, Welsh theorem: For any  $k > 0$  let  $A_{n,k}$  be the event that the box  $[0, kn] \times [0, n]$  contains a horizontal open crossing and let  $p \geq p_c$ . Then there exists a constant  $\delta_k$ , independent of  $n$  and  $p$  such that  $\mathbb{P}_p(A_{n,k}) > \delta_k$ .
- Consequence: For all  $p \geq p_c$ , the origin is surrounded by infinitely many open circuits with probability one.

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- Assign again *uniform* $[0,1]$  random variables to each edge independently, denoted by  $\tau(e)$  for an edge  $e$ .
- The invaded region  $\mathcal{S}(v)$  of a vertex  $v$  is defined as the increasing union of subgraphs  $\mathcal{S}_n(v)$ , where
  - $\mathcal{S}_0(v) = \{v\}$
  - $\mathcal{S}_{n+1}(v)$  is  $\mathcal{S}_n(v)$  together with the lowest edge not in  $\mathcal{S}_n(v)$  but incident to some vertex in  $\mathcal{S}_n$
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## Results for invasion percolation

- For all  $p > p_c$  we have that  $\mathcal{S} \cap \mathcal{C}_p \neq \emptyset$ , since the origin is surrounded by infinitely many  $p$ -open circuit.
- It is clear from the invasion mechanism that if  $\mathcal{S}_n \cap \mathcal{C}_p \neq \emptyset$  for some  $n > 0$  then  $\mathcal{S} \setminus \mathcal{S}_n \subset \mathcal{C}_p$
- By other words: If for any  $p$  the invasion hits the infinite  $p$ -open cluster, it will never leave this cluster again.
- If  $e_i$  is the edge invaded at time  $i$  then  $\limsup_{i \rightarrow \infty} \tau(e_i) = p_c$ .
- Since there is no percolation at  $p_c$  we get that  $\hat{\tau} = \max_{e \in E_\infty} \tau(e)$  exists and is greater than  $p_c$ .

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# Definition of the first pond

- Let  $\hat{e}$  be the edge where the maximum value of  $\tau$  is taken, namely  $\tau(\hat{e}) = \hat{\tau}$ .
- $\hat{e}$  exists and it is well-defined with probability 1.
- Suppose that  $\hat{e}$  is added to the invasion at time  $i + 1$ , then the graph  $\mathcal{S}_i = \hat{V}_1$  is called the first pond of the invasion or the first pond of the origin.

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# Further ponds

- Assume that the first pond is  $\mathcal{S}_{i_1}$  for some  $i_1$ .
- Then  $\max_{e_i \in E_\infty, i > i_1} \tau(e)$  exists and greater than  $p_c$ .
- Let  $\hat{e}_2$  be the edge where this value is taken.
- If  $\hat{e}_2$  is invaded at time  $i_2 + 1$ , then the graph  $\mathcal{S}_{i_2} \setminus \mathcal{S}_{i_1}$  is the second pond of the invasion.
- The other ponds of the invasion can be defined in similar way.

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